

MTH314: Discrete Mathematics for Engineers

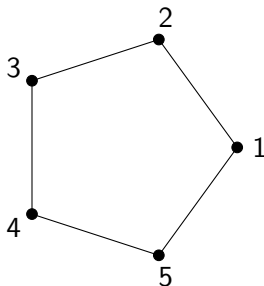
Lecture 9b: Introduction to Graph Theory

Dr Ewa Infeld

Ryerson University

Graph Theory Basics

A graph $G = (V, E)$ is a data structure/mathematical object that consists of a set of vertices/nodes V and a relation E (edges) on this set.



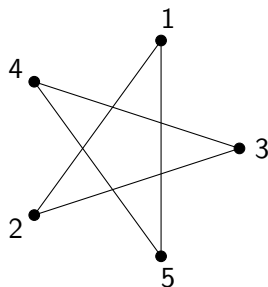
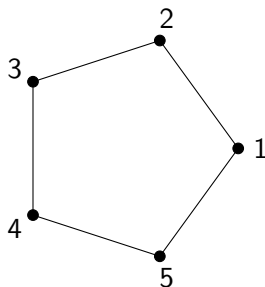
V is the vertex set

$$V = \{1, 2, 3, 4, 5\}$$

E is the set of edges

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These two are the same graph. Just two different ways to draw it.

Graph: Definition

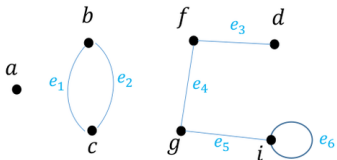
A graph $G = (V, E)$ is a data structure/mathematical object that consists of a set of vertices/nodes V and a relation E (edges) on this set.

Every edge $e \in E$ goes between two vertices, which we call *endpoints*. An edge from a vertex to itself is called a *loop*.

If there exists an $e \in E$ with endpoints $u, v \in V$ we say that u and v are adjacent, or that u is adjacent to v . We say that e is incident to both u and v .

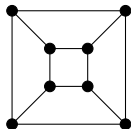
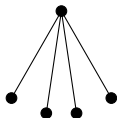
Example: $G = (V, E), V = \{a, b, c, d, f, g, i\}, E = \{e_1, e_2, e_3, e_4, e_5, e_6\}$

Edge	e_1	e_2	e_3	e_4	e_5	e_6
Endpoints	$\{b, c\}$	$\{b, c\}$	$\{f, d\}$	$\{f, g\}$	$\{g, i\}$	$\{i\}$

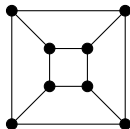
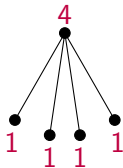


- e_6 is called a **loop**
- Vertex i is **adjacent** to itself
- Vertex a is **isolated**
- Edges e_1, e_2 are **parallel**
- Edges e_4, e_5 are **incident** to vertex g
- Vertices f, g are **connected** by edge e_4

The *degree* of a vertex is the number of edges coming out of that vertex. A loop will count twice, since both endpoints are at the same vertex.

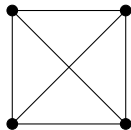


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A graph $G = (V, E)$ is a *simple* if there are no edges from a vertex to itself (“loops”) and between any two vertices there is at most one edge.

A *clique* is a simple graph where any two vertices are adjacent.

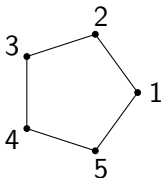


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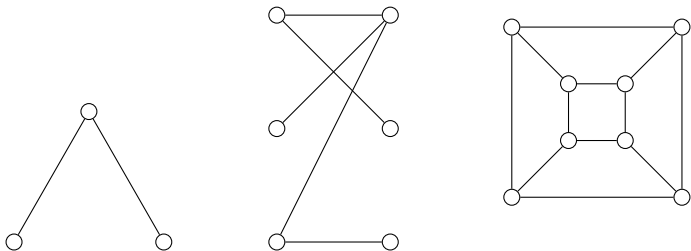


Not a clique:

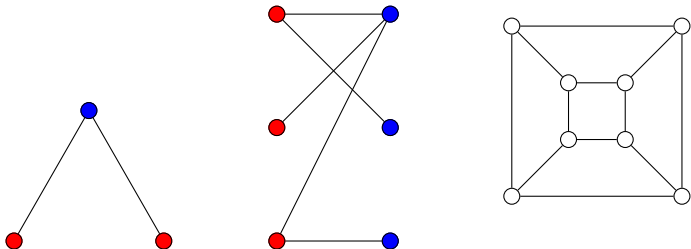


For example, vertices 1 and 3 are not adjacent.

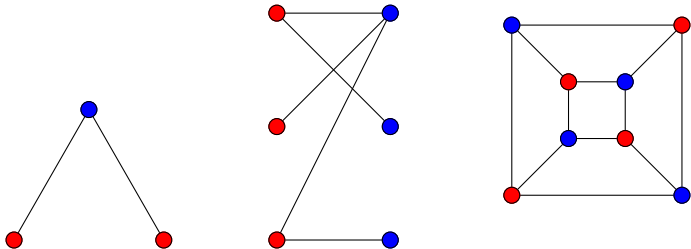
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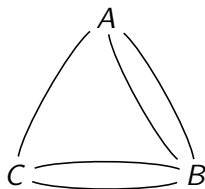


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Adjacency Matrix

An adjacency matrix is an integer matrix that encodes the graph. Rows correspond to vertices, and columns correspond to vertices. i, j -entry (i th row and j th column) is the integer representing how many edges connect vertices i and j .



$$\begin{array}{c} \text{A} \\ \text{B} \\ \text{C} \end{array} \begin{bmatrix} \text{A} & \text{B} & \text{C} \\ 0 & 2 & 1 \\ 2 & 0 & 2 \\ 1 & 2 & 0 \end{bmatrix}$$

In a simple graph, all entries are either 0 or 1 and all diagonal entries are 0. (why?)